Internet Connection Splitting: What's Old is New Again

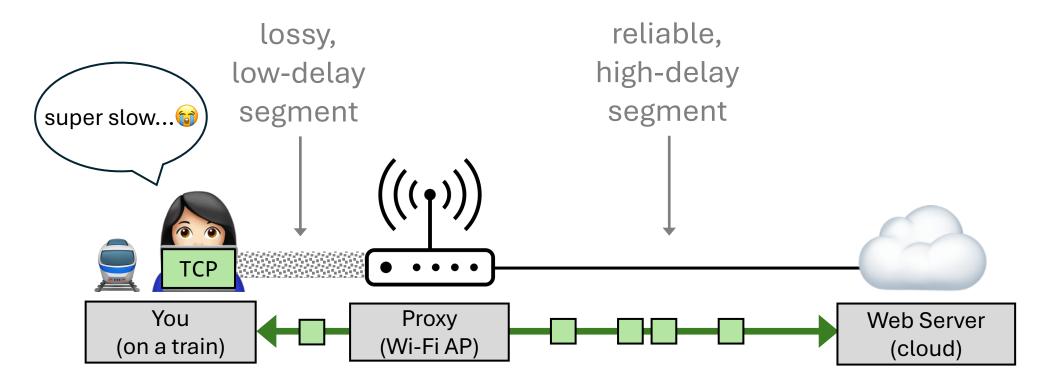
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Internet connection-splitting – E2E connection

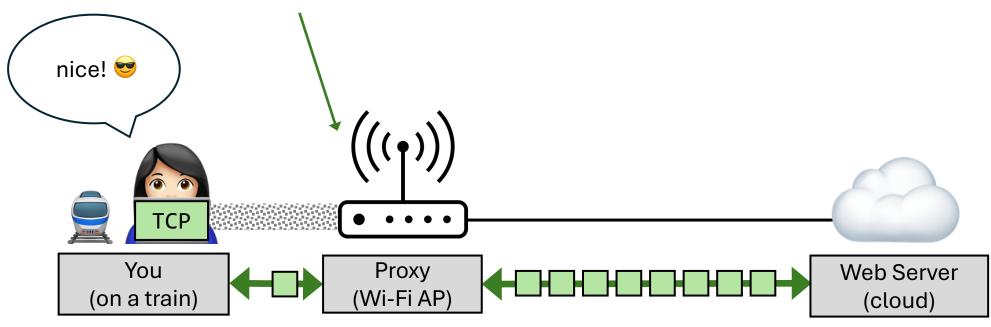


• Traditional "loss-based" congestion-control algorithms (e.g. CUBIC) treat wireless loss as network congestion in an end-to-end setting

Internet connection-splitting – split connection

connection-splitting TCP PEP

(performance-enhancing proxy)

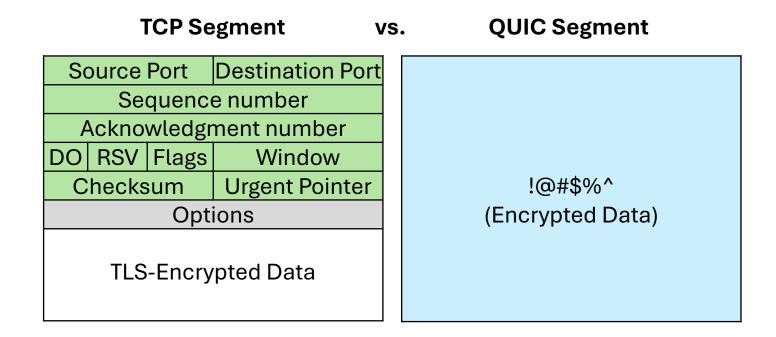


- Customized congestion control ("split CUBIC") on each path segment
- Widely deployed: 20-40% of paths, especially satellite and wireless^{1,2}

A lot has changed since these PEPs were first deployed in the 1990s...

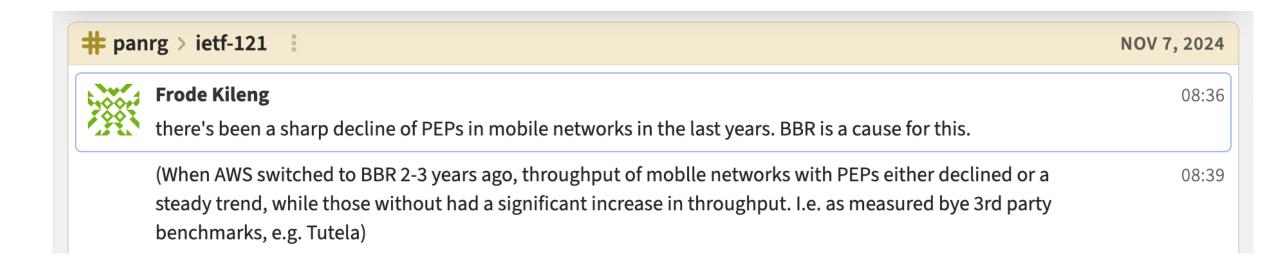
QUIC made connection-splitting impossible.

Encrypted transport protocol designed to allow the protocol to evolve over time and to improve performance for HTTPS traffic.



BBR made connection-splitting unnecessary.

Congestion-control algorithm that de-emphasizes loss as a congestion signal.



Is it true? Did BBR and QUIC make connection-splitting obsolete?

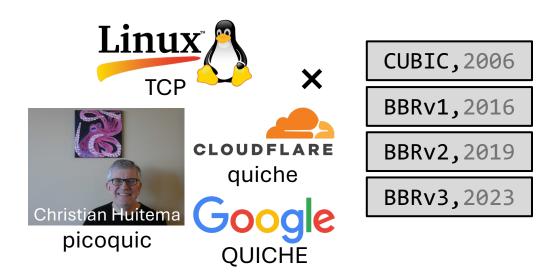
Re: Is splitting obsolete? It's complicated!

- Finding 1: TCP BBR has benefitted more from splitting over time.
- Finding 2: TCP BBRv3 even benefits in some classes of network paths where TCP CUBIC does not.
- Finding 3: QUIC implementations of the "same" congestion-control schemes vary significantly.

Measurement Study

- Plan: Evaluate various congestion-control schemes in a variety of network settings, both with and without a connection-splitting PEP.
- Question: Does the PEP improve performance with BBR/QUIC?

implementation/protocol/algorithm/version



Measurement Methodology

- 1. Pick a congestion control scheme.
- 2. Emulate two-segment topology: bandwidth, loss, delay.
- 3. Compare end-to-end vs. split throughput of HTTPS download.



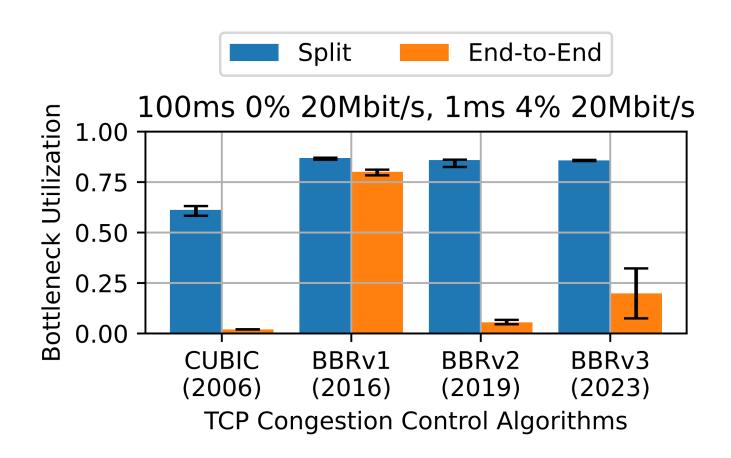
End-to-end connection

Split connection

Limitations

- 100% emulation study using tc-netem (not real-world)
- Metric is long-lived, full-throttle throughput only (not latency)
- Single-flow CCA environment (has fairness implications)

Finding 1: TCP BBR has benefitted more from splitting over time.

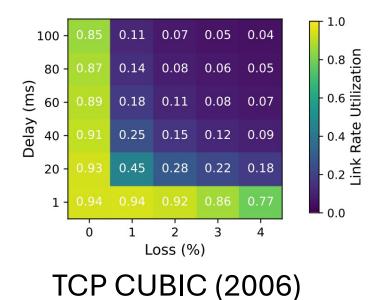


Why? BBRv2+ responded to concerns about TCP friendliness in BBRv1.

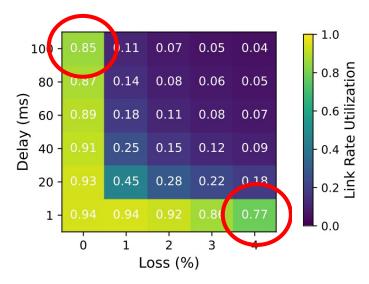
Does this generalize to other network settings?

Split Throughput Heuristic: Estimate *split throughput* as the <u>minimum</u> of the measured *end-to-end* throughputs on each segment of the split path.

Cache end-to-end measurements in a heatmap



Apply the heuristic to estimate split throughput

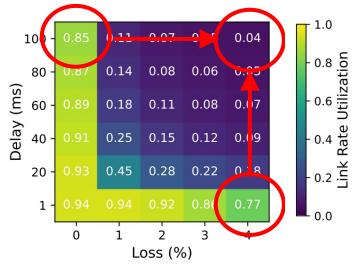


Split Utilization

min(0.77, 0.85) = 0.77

TCP CUBIC (2006)

Estimate end-to-end throughput in a split path



TCP CUBIC (2006)

Split Utilization

min(0.77, 0.85) = 0.77

End-to-End Utilization

bw = min(bw1, bw2) = min(10, 10) = 10 delay = delay1 + delay2 = 100 + 1 \approx 100 loss \approx loss1 + loss2 = 0 + 4 = 4

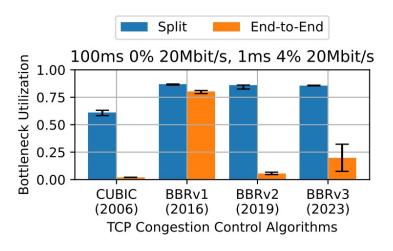
Finding 2: TCP BBRv3 even benefits in some classes of network paths where TCP CUBIC does not.

 $15 \cdot 21 \cdot 25 = 7875 \text{ split settings}$

 $5 \cdot 6 \cdot 5 = 150$ end-to-end settings

BBRv1	CUBIC	BBRv3
7875	7875	7875
0	2231	234
0	942	188
0	942	38
0	0	72
0	0	78
	7875 0 0	7875 7875 0 2231 0 942 0 942

Finding 2: TCP BBRv3 even benefits in some classes of network paths where TCP CUBIC does not.



Class I. Asymmetric, last-mile traditional PEP deployments with wireless link or rate policer



Split End-to-End

100ms 2% 20Mbit/s, 1ms 2% 20Mbit/s

1.00

0.75

0.50

0.25

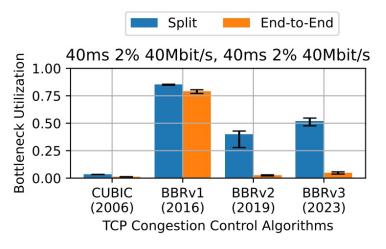
0.00

CUBIC BBRv1 BBRv2 BBRv3 (2006) (2016) (2019) (2023)

TCP Congestion Control Algorithms

Class II. Asymmetric, lossy low-resource networks, regions with no IXPs

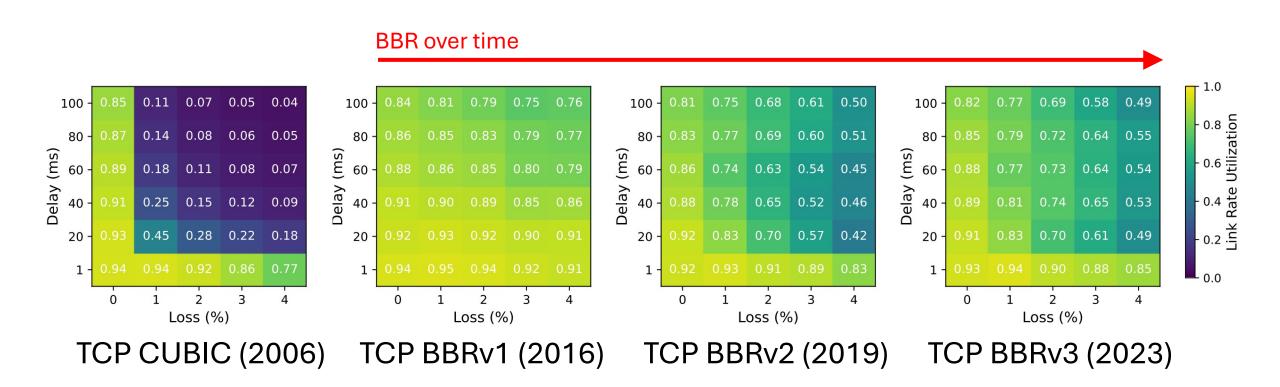




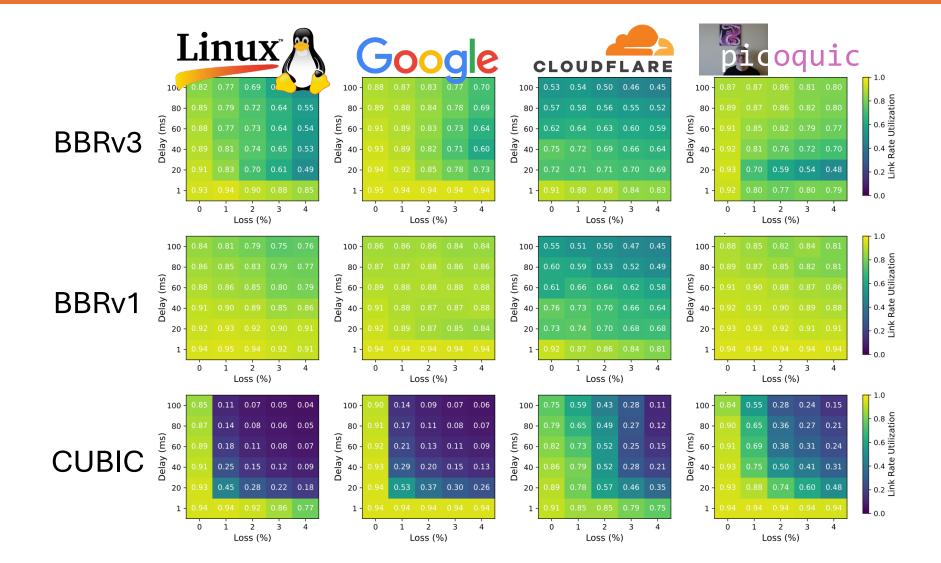
Class III. Symmetric, lossy wireless ad-hoc and satellites with lossy "middle-miles"



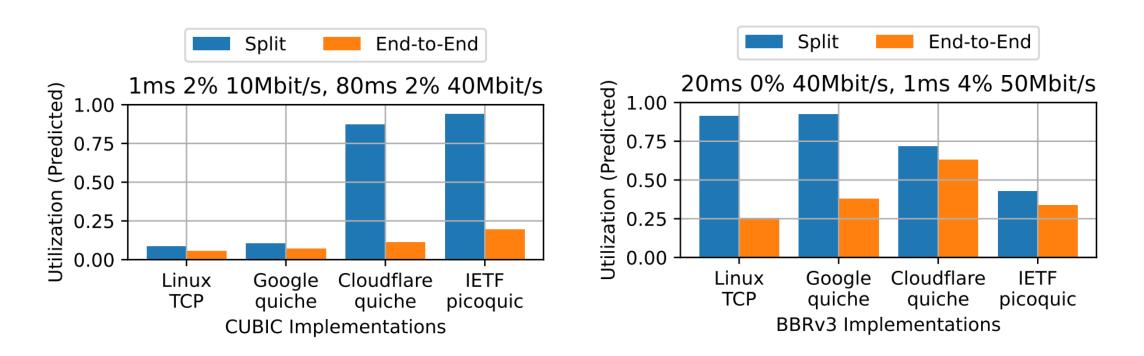
Heatmaps for characterizing end-to-end behavior



Finding 3: QUIC implementations of the "same" congestion-control schemes vary significantly, and with Linux TCP's.



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These results have implications for the throughput of the various user-space(!) QUIC implementations in *split* network path scenarios.

Calls to Action

- End-to-end congestion control
 - Refer to CC schemes by implementation/protocol/algorithm/version, e.g. not just "BBR" or even "QUIC BBRv1", but "Chromium QUIC BBRv1"
 - Standardize what it means to conform to a particular CCA standard, perhaps by creating performance test suites
- Connection-splitting
 - Protocol-agnostic ways to emulate PEPs⁵
 - Real-world studies to validate speculated performance improvements
 - Study other metrics improved by PEPs

Summary









- 1. TCP BBR has benefitted more from splitting over time since "v1" in 2016.
- 2. TCP BBRv3 even benefits in some settings where TCP CUBIC does not.
- 3. QUIC implementations of the same CC schemes vary significantly.

Calls to action: Address disparities between different implementations of the same congestion-control algorithms, and develop protocol-agnostic methods to emulate the performance benefits of PEPs.

https://github.com/StanfordSNR/connection-splitting